Data-driven systems

Query optimization



Contents

- Purpose of query optimization
- Microsoft SQL Server
 - > Execution plans
 - > Join and table access options
 - > General recommendation
- MongoDB
 - > Indices and execution plans

What is the purpose of query optimization?



What is the purpose of query optimization?









What is the purpose of query optimization?











Response time is affected by

- I/O cost
 - > Most prominent in data bases
 - > Does not improve according to Moore's law
 - > Needs special tricks
- CPU usage
 - > Complex queries
 - > Complex computations
- Memory usage
 - > Cache effect



Microsoft SQL Server

General concepts



Basics of the optimization

- Evaluates based on statistics
 - > Cost = response time (CPU + I/O time)
- Trivial plan
 - > Unambiguous for simple queries
 - > Rule-based
- When no trivial plan is available
 - > Complex queries
 - > Three phase optimization



Three phase optimization

- No trivial plan
- O. Phase
 - > Simple optimizations
 - > Preferred hash join
 - > If cost < X \rightarrow execute
- 1. Phase
 - > Complex optimizations
 - > If $cost < Y \rightarrow execute$
- 2. Phase
 - > Parallel execution



Process of executing a query

- Analyzer
 - > Compiler
 - > Logical plan
- Optimization
 - > Physical execution plan
 - > Read tables
 - > Joining tables
- Row executor
 - > Mapping physical plan to I/O operations
- Executor
 - > Executes the operations



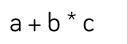
Microsoft SQL Server

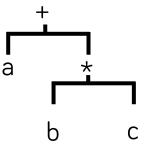
Logical execution plan



Elements of the logical execution plan

- Parser tree
 - > Relations (leaf)
 - > Operations (node)
 - > Data flow from bottom to top
- Relational algebra operations
 - > Cartesian-join (R×S)
 - > Projection $(\pi_1(R))$
 - > Selection $(\sigma_F(R))$
 - > Join (R⋈S)
 - > Filtering duplicates ($\delta(R)$)
 - > Grouping $(\gamma_L(R))$
 - > Sorting $(\tau_L(R))$





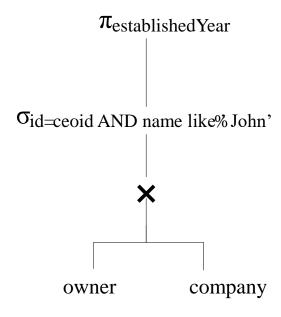


Parser tree

select establishedYear

from owner o, company c

where o.id = c.ceoid AND name like '%John'





Refactoring the parser tree - 1

- Create an optimal logical execution plan
- Reduce possible physical execution options
- Basic concept
 - > Move selection (where) down in the three
 - > Using joins
 - Cartesian join only when explicitly specified
 - > One side of a join should be a table



Refactoring the parser tree - 2

- Selection
 - > Can be re-ordered: $\sigma_{F1}(\sigma_{F2}(R)) = \sigma_{F2}(\sigma_{F1}(R))$
 - > Can be re-written:
 - $\sigma_{F \text{ and } G}(R) = \sigma_{F}(\sigma_{G}(R))$
 - $-\sigma_{F \text{ or } G}(R) = \sigma_{F}(R) \text{ UNION } \sigma_{G}(R)$
- Join
 - $> R\bowtie_F S = \sigma_F(R \times S)$
 - $> R \bowtie S = S \bowtie R$
 - $> (R \bowtie S) \bowtie U = R \bowtie (S \bowtie U)$

Refactoring the parser tree - 3

Joining and selection

$$> \sigma_F(R \bowtie S) = \sigma_F(R) \bowtie S$$

- If R has all attributes of F

$$> \sigma_{F}(R \bowtie S) = R \bowtie \sigma_{F}(S)$$

- If S has all attributes of F

$$> \sigma_F(R \bowtie S) = \sigma_F(R) \bowtie \sigma_F(S)$$

- If both R and S has attributes of F
- Duplicates

$$> \delta(\gamma_L(R)) = \gamma_L(R)$$

$$> \delta(R \times S) = \delta(R) \times \delta(S)$$
 (similarly for joins)

Microsoft SQL Server

Physical execution plan



Physical plan

- Elements of the physical plan
 - > Operators for reading the table
 - Logical plan leaf reading operations
 - > Executing relational algebra operations
- Creating plans
 - > Rule-based
 - > Cost-based
 - Table seek methods
 - Implementation of joins
 - Order of joins



Nested loop join

- Two embedded for cycles
- I/O cost
 - > O(num_block_1 * num_block_2)
- Works in all cases
 - > In case of large tables: keep only partitions of the tables in memory



Hash join

- First pass
 - > Read the smaller table
 - > Build a hash table in memory
 - Key is the column used for joining
- Second pass
 - > Read the larger table
 - > Search for matching records in the hash table
- I/O cost
 - > O(num_block_1 + num_block_2)



Sort Merge Join

- Reads both tables into memory
- Sorts based on the joined columns
- Merge the two sorted lists
 - > While "walking" the two sorted lists
- For small tables
- Index due to sorting
- I/O cost
 - > O(num_block_1 + num_block_2)

Table scan methods - 1

- Generally there are two ways
 - > Full scan
 - For small tables
 - When most rows are needed
 - > Index scan
 - When using filtering
 - If there is an index covering the filter criteria
 - Sorting



Table scan methods - 2

- Table scan
 - > No index
 - > Evaluates the filtering condition
- Clustered index scan
 - > Clustered reading
 - > Data blocks ordered by index
 - > Clustered index created along primary key
 - > Preferred over table scan



Table scan methods - 3

- Nonclustered index scan
 - > Similar to clustered index scan
 - > Mostly for evaluating =
- Clustered/Nonclustered index seek
 - > Similar to index scan
 - > Walks the index from a starting point
 - ->, between, < operators



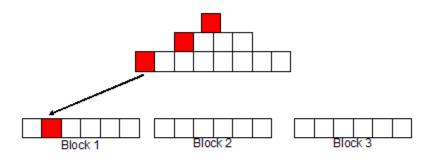
MS SQL Server indexes

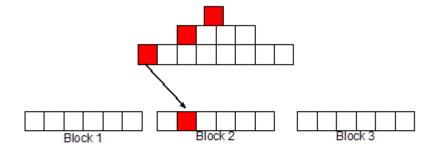
- B* tree
 - > Simple
 - > Compound indexes
 - Hierarchical
 - > Clustered
 - Data blocks orderes by index
 - One for each table
 - Automatically created for the primary key



MS SQL Server indexes

Clustered / non-clustered





Select order_nbr, item_name from ordor natural join item;

Clustered table rows

Clustering_factor ~= blocks

Select order_nbr, item_name from ordor natural join item;

Un-Clustered table rows

Clustering_factor ~= num_rows

Source: http://www.dba-oracle.com/t_table_row_resequencing.htm



Indexes - 1

- Cover index (included column)
 - > Adding further data into the B* tree leaf nodes
 - > The row data does not need to be accessed
- Using clustered and non clustered indices together
 - > Nonclustered index leaf
 - Does not contain a physical address
 - Points to the clustered index
 - > Double index read



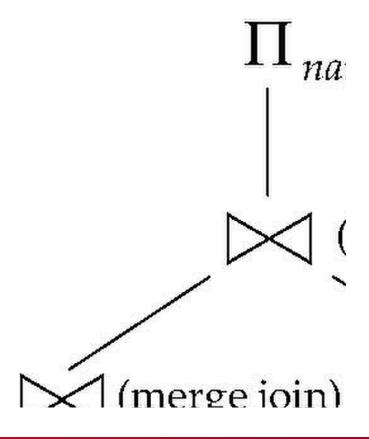
Indexes - 2

- Indexed views
 - > The view result is stored
 - > Index can be defined, works as for tables



Execution plan

 It defines exactly what action is performed on each node





Végrehajtási terv megnézése

Scan Direction

```
■ New Query ■ 🔊 🛣 🛣 🛣 🛣 🗂 🗂
select p.name from Product p
                                                                 ▶ Execute ■ ✓ 등 🗊 🖫 등 등
join Category c on p.CategoryID = c.ID
where c.Name = 'LEGO'
                                                               Include Actual Execution Plan (Ctrl+M) pe (54))*
                   \pi_{\mathsf{name}}

    ■ Results    ■ Messages    □ Execution plan

                                                             Query 1: Query cost (relative to the ba
                                                             select p.name from Product p join Cated
                                                                         Nested Loops
                                                                                        Clustered Inde...
                                                                                        [Product].[PK ...
                                                                         (Inner Join)
                                                                          Cost: 1 %
                                                                                          Cost: 41 %
                                                             Cost: 0 %
                                                                           0.000s
                                                                                            0.000s
                                                                            1 of
                                                                                            10 of
                                                                           1 (100%)
                                                                                          10 (100%)
        Oname='LEGO'
                                                                                        Clustered Inde…
                                                                                        [Category].[PK...
                                                                                          Cost: 58 %
                              Product
           Category
                                                                                            0.000s
                                                                                           10 (10%)
                    Parallel
                                        False
                    Physical Operation
                                        Clustered Index Seek
                    Predicate
                                        [M22XDS].[dbo].[Category].[Name] as [c].[Name]=N'LEGO'
```

FORWARD

Query plan in SQL Server

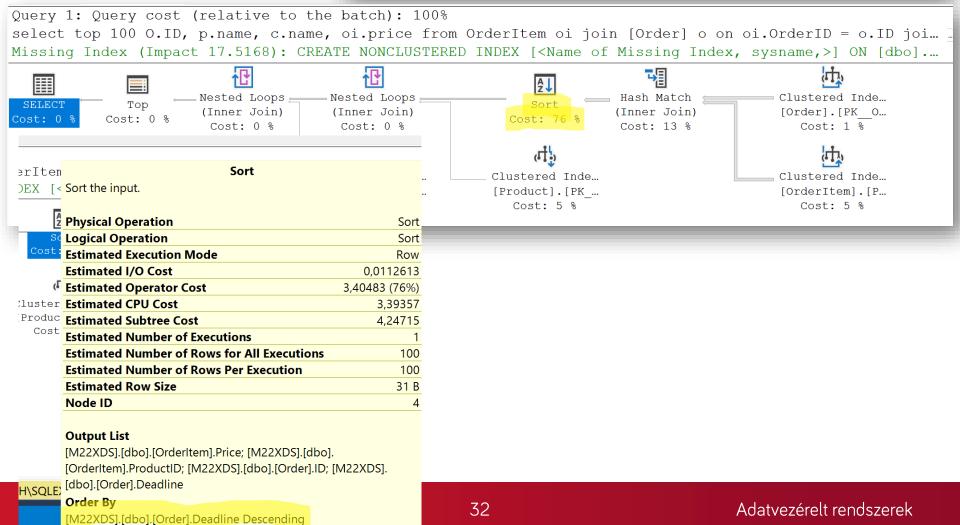
Press CTRL+L after executing the query

```
from
   OrderItem oi join [Order] o on oi.OrderID = o.ID
   join Product p on oi.ProductID = p.ID
   join Category c on p.CategoryID = c.ID
```

```
Query 1: Query cost (relative to the batch): 100%
select top 100 O.ID, p.name, c.name, oi.price from OrderItem oi join [Order] o on oi.Ord
                         Nested Loops
                                         — Nested Loops ====
                                                                  Nested Loops
                                                                                        Clustered Inde...
               qoT
                         (Inner Join)
                                           (Inner Join)
                                                                  (Inner Join)
                                                                                        [OrderItem].[P...
            Cost: 0 %
                           Cost: 0 %
                                            Cost: 0 %
                                                                    Cost: 0 %
                                                                                           Cost: 1 %
                                         Clustered Inde...
                                                                 Clustered Inde...
                                                                                        Clustered Inde...
                                          [Category].[PK...
                                                                 [Product].[PK ...
                                                                                        [Order].[PK O...
                                            Cost: 5 %
                                                                   Cost: 57 %
                                                                                          Cost: 36 %
```



Index hints

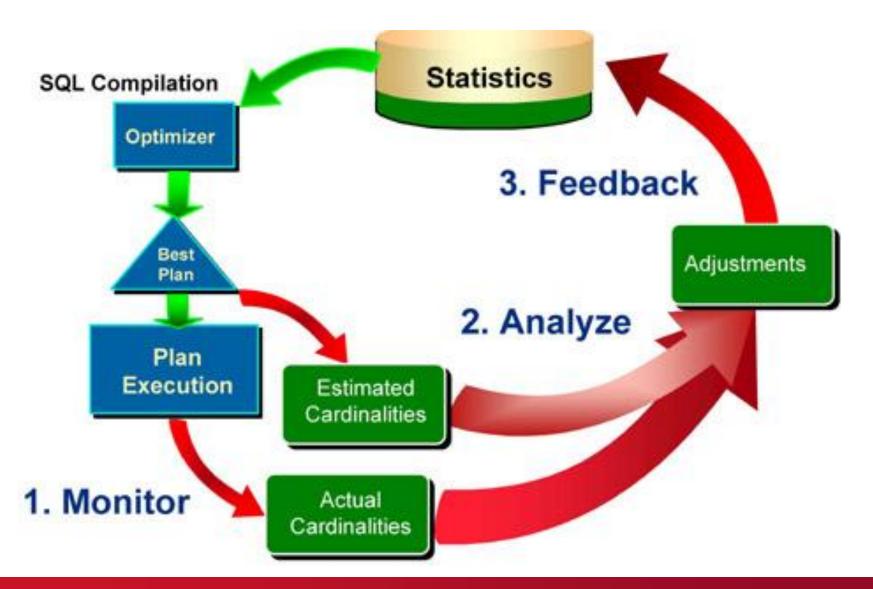


Execution plan alternatives

- There can be several plan alternatives, which one is optimal?
 - > Huge differences: seconds or days
 - > The cost depends on how many lines are the result for each phase
 - > The system estimates this based on statistics
 - > Self-tuning database, redesigns if necessary, during execution
- Plan cache
 - > Execution plan cache
 - > Used when for queries with the same structure
 - > Statistics have not changed
 - You may need to update the cache manually!
 - The statistics maintenance must be setup!



Self-tuning





Microsoft SQL Server

General recommendation



Best practices - 1

- Keep statistics up to date
 - > Deprecated statistics -> suboptimal execution plan
 - > (This is the default unless turned off)
- Structure of the query
 - > SQL is declarative
 - Keep procedural execution in mind
 - > Same result can be obtained multiple ways
 - > Make it simple
 - > Avoid select *
 - > Good structure is advantageous
 - Use hints as a last resort



- Prefer join over
 - > In / Not in
 - > Exists / Not exists
- Prefer In over Exists
- Views
 - > Avoid if possible
 - > Do not join them
- Avoid Or clauses → Union all
- Union all if possible



```
select *
from Invoice i
where not exists
(
select 1
from InvoiceItem ii
where i.Id=ii.InvoiceID
)
```

```
select i.*
from Invoice i
where i.id not in
(
select InvoiceID
from InvoiceItem
)
```

```
select i.*
from Invoice i left outer join InvoiceItem ii
on i.Id=ii.InvoiceID
where ii.id is null
```

Does not matter in simple cases



- Using indexes
 - Usually one can be used for one table in a query
 Join may use it up
 - > Compound indexes
 - Hierarchy matters
 - > If the key is used in an expression, the optimizer cannot use it
 - E.g., key+0 (\leftarrow Optimizer may handle this though)



- Using functions
 - > No problem in a select
 - Does not affect the execution plan
 - > Avoid in a where clause
 - Has to be evaluated for each record
 - Hard to move in the query tree
 - No statistics are available for its output → hard to optimize



Further reading material

- Grant Fritchey: SQL Server Execution Plans, Simple Talk Publishing, 2012
 - > pdf: http://www.red-gate.com/community/books/sql-server-execution-plans-ed-2



MongoDB

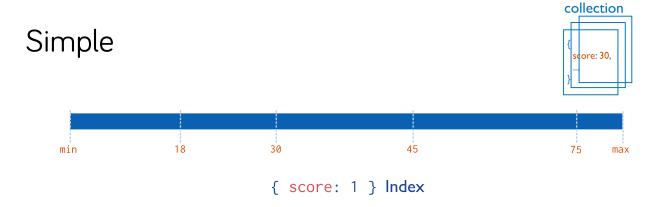


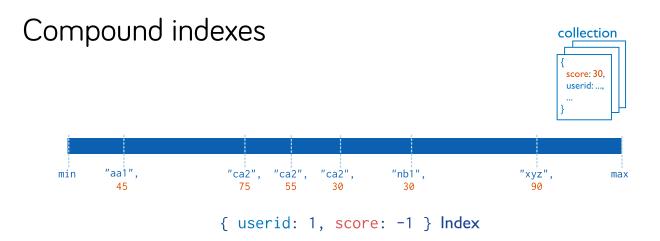
Indices in MongoDB

- Index "only" for lookup
 - > (As there is no join operation)
- Types of indices
 - > Simple and compound
 - > Unique index
 - Can be used to ensure a primary key-like attribute
 - > Indexes content of arrays too
 - > Indexes nested objects too
 - > TTL, Geospatial, full text
- Index must be defined
 - > Except: _id unique



Types of indices





Images source: https://docs.mongodb.com/manual/indexes/

Basics of the optimization

- Does not use statistics
- Choice between multiple possible plans
 - > Starts execution all, whichever yields the first 101 results, is the best
- How can there be multiple plans?
 - > There are multiple indices covering the query



Optimization steps

- Move filtering ahead of other steps
 - > Before projection, and if needed, split the filtering into two
 - > Before sorting
- Move skip and limit ahead
 - > Projection
- Merge
 - > Limit + limit, skip + skip

Plan cache

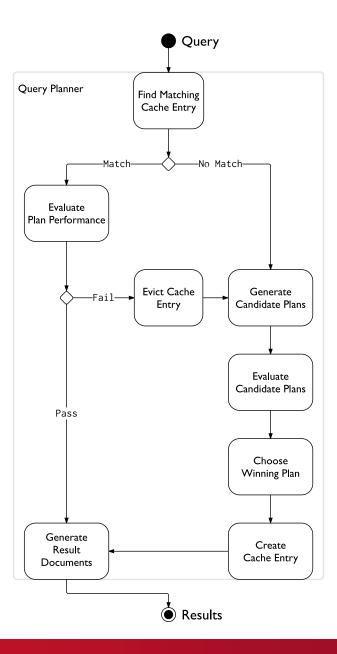


Image source:

https://docs.mongodb.com/manual/core/query-plans/



Plan cache

- Plan cache
 - > Structurally similar plans
 - > Pass/fail evaluation
- Query shape
 - > Filters, sorting, etc. used
 - > No values
 - E.g. for filtering only the filtered field name is present



Explain

• query.explain()

```
"winningPlan" : {
   "stage" : <STAGE1>,
   "inputStage" : {
      "stage" : <STAGE2>,
      "inputStage" : {
         "stage" : <STAGE3>,
"rejectedPlans" : [
   <candidate plan 1>,
```

Stage-ek

- COLLSCAN
- IXSCAN
- FETCH

• ...

